

Synthesis of Recursive Functions from Sequences of Natural Numbers¹

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Problem

Given a “finite” sequence of natural numbers,

A000217 : 0, 1, 3, 6, 10, 15, 21, 28, 36, 45, 55, 66, 78, 91, 105, 120

generate a “small” program that match the sequence?

$$f(x) = \frac{x \times (x + 1)}{2}$$

Motivations

1) Conjecturing:

$$\sum_{x=0}^n x = \frac{x \times (x + 1)}{2} \quad ?$$

2) A shorter explanation generalizes better.

3) A shorter explanation gives some understanding.

$$\sum_{x=0}^n x \text{ is not prime for } x > 2$$

Language: understandable, efficient, generalizes

one variable: x

constants: 0, 1, 2

functions: $+$, $-$, \times , $/$, *mod*, $\sqrt{\quad}$, *power*

conditional statements:

- $\text{cond}(a, b, c) = \text{if } a = 0 \text{ then } b \text{ else } c$

- $\text{loop}(f, a, b) = f^a(b)$

- $\text{halt}(f, a) = \text{minimum } i \text{ such that } f^i(a) = 0$

A issue with linear synthesis?

Same sub-expression repeatedly synthesized.

$$f(x) \times (x \times (x + 1)) \times g(x)$$

$$(x \times (x + 1)) + h(x)$$

Factorized bottom-up synthesis

target T : $[0, 1, 3, 6, 10, 15, \dots]$

size 1: $x, 0, 1, 2$

size 2: $\sqrt{x}, \sqrt{0}, \sqrt{1}, \sqrt{2}$

size 3: $x + x, x + 1, 2 \times x, \dots$

size 4: $loop(\lambda x. x, 1, 2), \sqrt{x + x}, \dots$

solution: program f such as $[f(0), f(1), \dots, f(15)] = T$

random fixed width $w \in \{4, 8, 16, 32\}$ for each search:
select w programs at each size.

Selecting programs for a target

target: $[0, 1, 3, 6, 10, 15, \dots]$

sub-program: $x + 1 \equiv [1, 2, 3, 4, 5, 6 \dots]$

$[1, 2, 3, 4, 5, 6 \dots]$ useful for $[0, 1, 3, 6, 10, 15, \dots]$?

Euclidean distance is not good

because $[2, 14, 3, \dots]$ is useful for $[2^2, 2^{14}, 2^3, \dots]$.

Train a classifier from solutions

Let P be a minimal solution and P_{pos} a subprogram of P .

positive example: $([P], [P_{pos}])$

Let P_{neg} be a generated program with the same size as P_{pos} which is not a subprogram of P .

negative example: $([P], [P_{neg}])$

OEIS sequences and restrictions

- 1) At least 16 elements.
- 2) First 16 elements between 0 and $2^{63} - 1$.
- 3) First 8 elements different from every other OEIS sequence.

About 350 000 sequences become about 200 000 targets.

Reinforcement learning

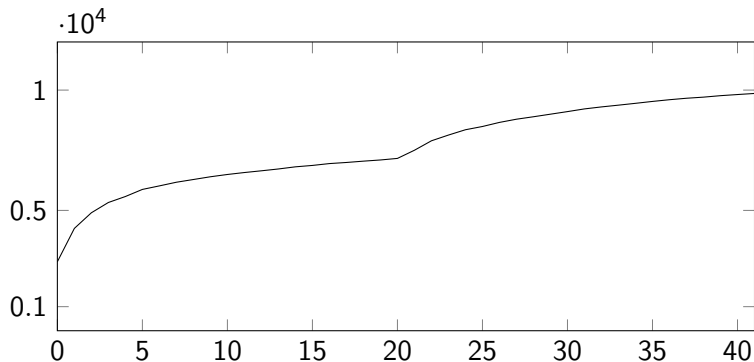


Figure: Number y of training problems solved after generation x

Solutions

Picked at random: A079273

Octo numbers

1, 10, 29, 58, 97, 146, 205, 274, 353, 442, ...

$$(x \times x) + (1 + (x + x))^2$$

Smallest with a nested loop: A125833

Numbers whose base 5 representation is 333.....3

0, 3, 18, 93, 468, 2343, 11718, 58593, 292968, ...

Def: $f(x) = 1 + (x + x)$, $g(x) = x + f^2(x)$

$$g^x(0)$$

Solutions

Largest: A273848

Number of active (ON,black) cells at stage $2^n - 1$ of the two-dimensional cellular automaton defined by "Rule 969", based on the 5-celled von Neumann neighborhood.

1, 4, 45, 225, 961, 3969, 16129, 65025, 261121, ...

Def: $f(x) = x + (1 + x)$

(if $x/2 = 0$ then 2^x else if $\sqrt{x+1}/2 = 0$ then 3^x else $f^x(1)$)

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(if $x/2 = 0$ then 2^x else if $\sqrt{x+1}/2 = 0$ then $f(x)$ else $f^x(1)$)

Conclusion

Factorized bottom-up program synthesis

Semantic quotient + semantic filtering (semantic = sequences)

Interesting results (requires more learning)

Future work

More extensive experiments:

- mix syntactic and semantic features
- large integers, real numbers, multiple variables, lists
- **backward reasoning, recursion**, techniques from ILP

Apply synthesis to theorem proving:

- term synthesis, tactic synthesis, cut introduction.

Look for applications beyond mathematics.