

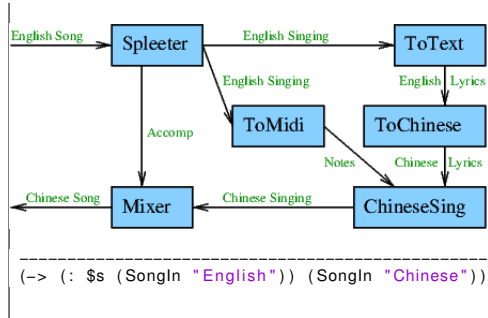
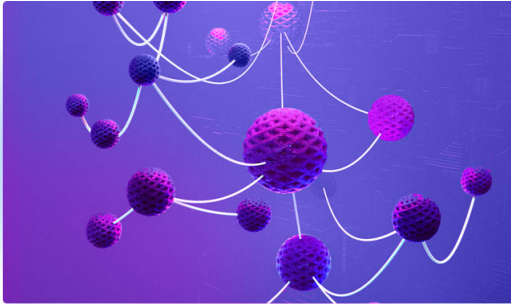
Meta-Reasoning in MeTTa

for Inference Control via Provably Pruning the Search Tree

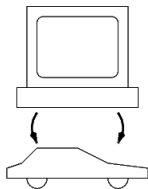
Nil Geisweiller

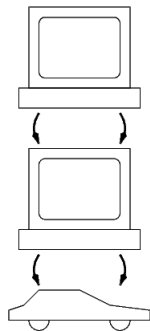


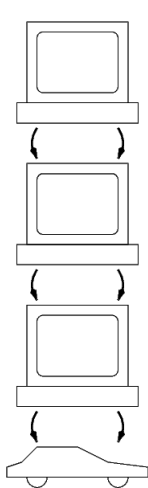
Artificial Intelligence and Theorem Proving 2024 (AITP-24)

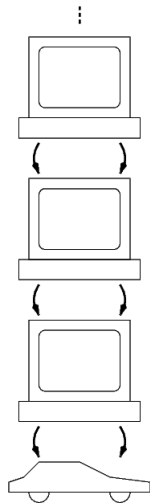


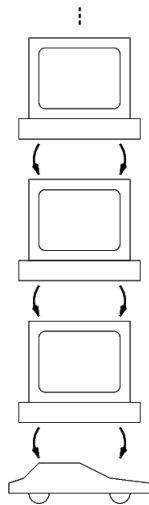
- MeTTa
- Meta-reasoning



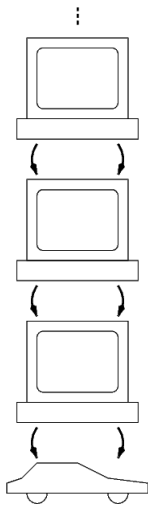






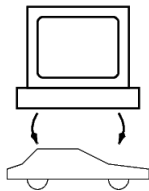


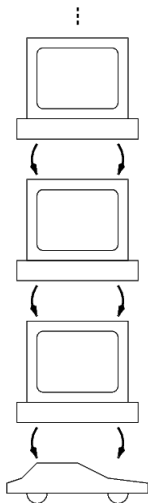
Gödel Machine (*Jürgen Schmidhuber, 2003*)



Gödel Machine (*Jürgen Schmidhuber, 2003*)

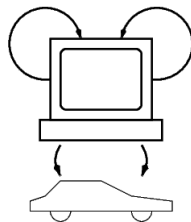
- Merge all machines into one.

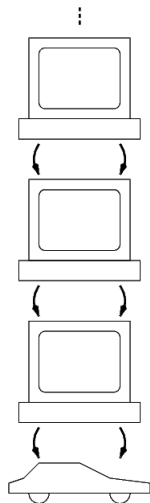




Gödel Machine (*Jürgen Schmidhuber, 2003*)

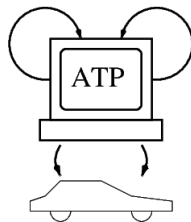
- Merge all machines into one.
- Internal actions to action space.





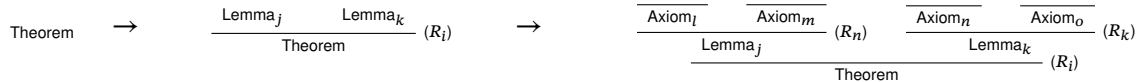
Gödel Machine (*Jürgen Schmidhuber, 2003*)

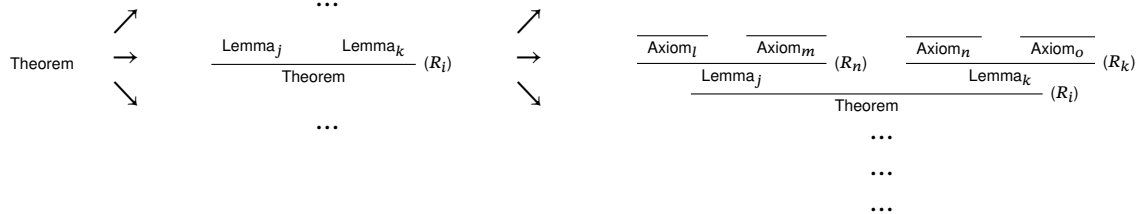
- Merge all machines into one.
- Internal actions to action space.
- Mathematical proof \Rightarrow trigger action.

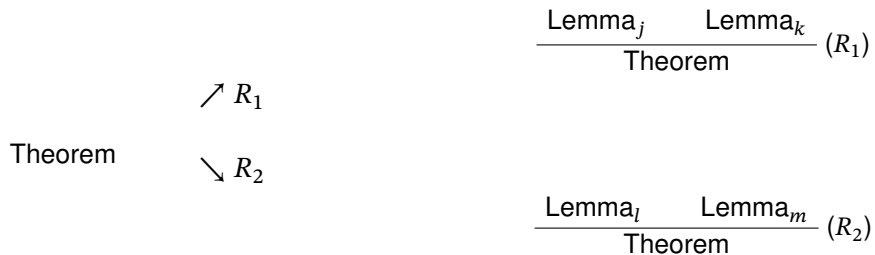


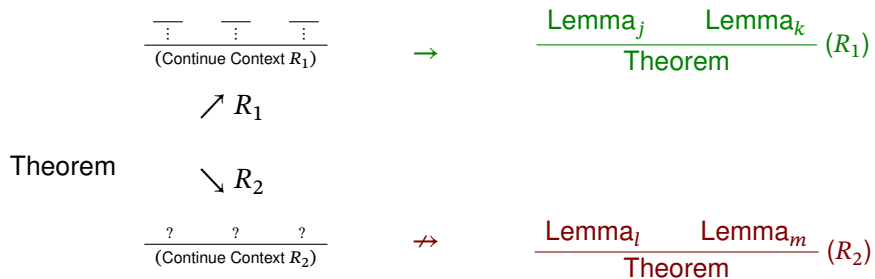
Theorem

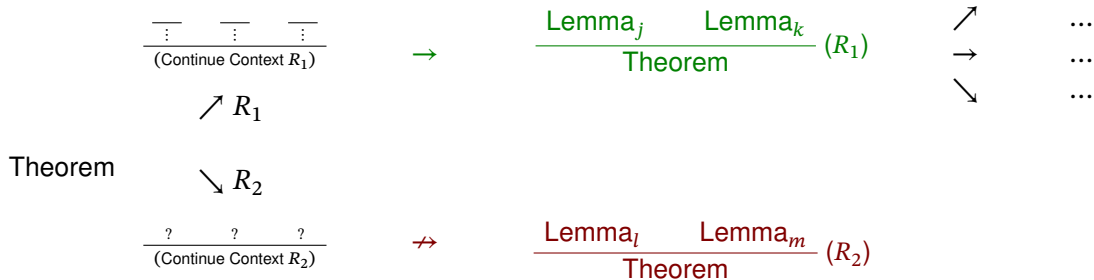
Theorem \rightarrow $\frac{\text{Lemma}_j \quad \text{Lemma}_k}{\text{Theorem}} (R_i)$











MeTTa: Meta Type Talk

- Functional and logic programming
- Non-determinism (like Curry)
- Unification (like Prolog)
- Gradual typing
- Self-modifiable
- Concurrency
- Scalable

```
;; Bit strings
(= (bits Z) Nil)
(= (bits (S $k)) (Cons 0 (bits $k)))
(= (bits (S $k)) (Cons 1 (bits $k)))

;; Generate all 3-bit strings
!(bits (S (S (S Z))))
```



```
[(Cons 0 (Cons 0 (Cons 0 Nil)))
 (Cons 0 (Cons 0 (Cons 1 Nil)))
 (Cons 0 (Cons 1 (Cons 0 Nil)))
 (Cons 0 (Cons 1 (Cons 1 Nil)))
 (Cons 1 (Cons 0 (Cons 0 Nil)))
 (Cons 1 (Cons 0 (Cons 1 Nil)))
 (Cons 1 (Cons 1 (Cons 0 Nil)))
 (Cons 1 (Cons 1 (Cons 1 Nil)))]
```

```
;; Backward chainer
(= (bc $kb $ _ (: $prf $ccln)) (match $kb (: $prf $ccln) (: $prf $ccln)))
(= (bc $kb (S $k) (: ($prfabs $prfarg) $ccln))
  (let* (((: $prfabs (-> $prms $ccln)) (bc $kb $k (: $prfabs (-> $prms $ccln))))
         ((: $prfarg $prms) (bc $kb $k (: $prfarg $prms))))
    (: ($prfabs $prfarg) $ccln)))
```

```
;; Backward chainer
(= (bc $kb $ _ (: $prf $ccln)) (match $kb (: $prf $ccln) (: $prf $ccln)))
(= (bc $kb (S $k) (: ($prfabs $prfarg) $ccln))
  (let* (((: $prfabs (-> $prms $ccln)) (bc $kb $k (: $prfabs (-> $prms $ccln))))
         ((: $prfarg $prms) (bc $kb $k (: $prfarg $prms))))
    (: ($prfabs $prfarg) $ccln)))

;; Knowledge base
!(bind! &kb (new-space))
!(add-atom &kb (: AK (-> $a (-> $b $a))))
!(add-atom &kb (: AS (-> (-> $a (-> $b $c)) (-> (-> $a $b) (-> $a $c)))))
```

```
;; Backward chainer
(= (bc $kb $ _ (: $prf $ccln)) (match $kb (: $prf $ccln) (: $prf $ccln)))
(= (bc $kb (S $k) (: ($prfabs $prfarg) $ccln))
  (let* (((: $prfabs (-> $prms $ccln)) (bc $kb $k (: $prfabs (-> $prms $ccln))))
        ((: $prfarg $prms) (bc $kb $k (: $prfarg $prms))))
    (: ($prfabs $prfarg) $ccln)))

;; Knowledge base
!(bind! &kb (new-space))
!(add-atom &kb (: AK (-> $a (-> $b $a))))
!(add-atom &kb (: AS (-> (-> $a (-> $b $c)) (-> (-> $a $b) (-> $a $c)))))

;; Query
!(bc &kb (S (S Z)) (: $prf (-> $a $a)))
```

```

;; Backward chainer
(= (bc $kb $ _ (: $prf $ccln)) (match $kb (: $prf $ccln) (: $prf $ccln)))
(= (bc $kb (S $k) (: ($prfabs $prfarg) $ccln))
  (let* (((: $prfabs (-> $prms $ccln)) (bc $kb $k (: $prfabs (-> $prms $ccln))))
         ((: $prfarg $prms) (bc $kb $k (: $prfarg $prms))))
    (: ($prfabs $prfarg) $ccln)))

;; Knowledge base
!(bind! &kb (new-space))
!(add-atom &kb (: AK (-> $a (-> $b $a))))
!(add-atom &kb (: AS (-> (-> $a (-> $b $c)) (-> (-> $a $b) (-> $a $c)))))

;; Query
!(bc &kb (S (S Z)) (: $prf (-> $a $a)))

;; Results
[(: ((AS AK) AK) (-> $a $a))
 ...]

```

```

;; Backward chainer with dependent types and lambda abstraction

;;;;;;;;;;;;;;;;;;;;;;;;
;; Base cases ;;
;;;;;;;;;;;;;;;;;;;;;;;;

;; Match the knowledge base
(= (bc $kb $env $idx $_ (: $prf $thrm))
   (match $kb (: $prf $thrm) (: $prf $thrm)))

;; Match the environment
(= (bc $kb $env $idx $_ (: $prf $thrm))
   (match' $env (: $prf $thrm) (: $prf $thrm)))

;;;;;;;;;;;;;;;;;;;;;;;;
;; Recursive steps ;;
;;;;;;;;;;;;;;;;;;;;;;;;

;; Proof application
(= (bc $kb $env $idx (S $k) (: ($prfabs (: $prfarg $prms)) $thrm))
   (let* (((: $prfabs (-> (: $prfarg $prms) $thrm))
          (bc $kb $env $idx $k (: $prfabs (-> (: $prfarg $prms) $thrm))))
         ((: $prfarg $prms)
          (bc $kb $env $idx $k (: $prfarg $prms))))
     (: ($prfabs (: $prfarg $prms)) $thrm)))

;; Proof abstraction
(= (bc $kb $env $idx (S $k) (: (\ $idx $prfbdy) (-> (: $idx $prms) $thrm)))
   (let (: $prfbdy $thrm)
     (bc $kb (Cons (: $idx $prms) $env) (s $idx) $k (: $prfbdy $thrm))
      (: ( $idx $prfbdy) (-> (: $idx $prms) $thrm))))

```



```
;; Equality is transitive
```

```
!(add-atom &kb (: Trans (-> (: $prf1 (=== $x $y)) ; Premise 1
                             (-> (: $prf2 (=== $y $z)) ; Premise 2
                                 (=== $x $z)))))) ; Conclusion
```

```
;; Equality is symmetric
```

```
!(add-atom &kb (: Sym (-> (: $prf (=== $x $y)) ; Premise
                             (=== $y $x)))) ; Conclusion
```

```
;; Equality respects function application
```

```
!(add-atom &kb (: Cong (-> (: $f (-> (: $_ $a) $b)) ; Premise 1
                             (-> (: $x $a) ; Premise 2
                                 (-> (: $x' $a) ; Premise 3
                                     (-> (: $prf (=== $x $x') ; Premise 4
                                         (=== ($f $x) ($f $x')))))))) ; Conclusion
```

```
;; Rule of replacement
```

```
!(add-atom &kb (: Replace (-> (: $prf1 (=== $x $x') ; Premise 1
                                (-> (: $prf2 $x ; Premise 2
                                    $x')) ; Conclusion
```

```
;; Define double
```

```
!(add-atom &kb (: double (-> (: $k N) N)))
!(add-atom &kb (: double_base (=== (double (: Z N)) Z)))
!(add-atom &kb (: double_rec (-> (: $k N)
                                (=== (double (: (S (: $k N)) N)) (S (: (S (: (double (: $k N)) N)) N))))))
```

```
...
```

```
;; Query: find proof that for any natural k, (double k) is even
!(bc &kb Nil z (fromNumber 11) (: $prf (-> (: $k N) (Even (double (: $k N))))))
```

```
;; Results
```

```
[(: ((SIN ((Replace (((Cong Even) Z) (double Z)) (Sym double_base))))))
  (\ z (\ (s z) ((Replace (Sym (((Cong Even) (double (S z))) (S (S (double z)))) (double_rec z))))
    (MkEvenSS (s z))))))
  (-> (: $k N) (Even (double (: $k N))))
  ...]
```

```

                                     -(z)
                                     N
                                     -(double)
                                     N
                                     -(S)
                                     N
                                     N
                                     -(z)
                                     N
-----(Even) -----(double) -----(S) -----(double_rec)
(-> (: $k N) Type) N N N (== (double (S z)) (S (S (double z)))) (Cong) -----(s z)
                                     (== (Even (double (S $k))) (Even (S (S (double $k))))) (Sym) (Even (double z))
                                     (== (Even (S (S (double $k)))) (Even (double (S $k)))) (Even (S (S (double z)))) (MkEvenSS)
-----(Even) -(Z) -----(double_base)
(-> N Type) N N N (== (double Z) Z) (Sym) (Even (double z)) (Replace)
-----(Even) -----(double) -----(S)
(-> N Type) N N N (== Z (double Z)) (Cong) -----(MkEvenZ) -----(s z) -----(double (S z))
(== (Even Z) (Even (double Z))) (Even Z) (Even (double z)) (Even (double (S z))) (\)
-----(Even (double Z)) -----(Replace) N (-> (: (s z) (Even (double z))) (Even (double (S z)))) (\)
                                     (-> (: z N) (-> (: (s z) (Even (double z))) (Even (double (S z)))))
-----(Even (double Z)) -----(SIN)
                                     (-> (: $k N) (Even (double $k)))

```

```
;; Backward chainer
(= (bc $kb $ _ (: $prf $ccln)) (match $kb (: $prf $ccln) (: $prf $ccln)))
(= (bc $kb (S $k) (: ($prfabs $prfarg) $ccln))
  (let* (((: $prfabs (-> $prms $ccln)) (bc $kb $k (: $prfabs (-> $prms $ccln))))
        ((: $prfarg $prms) (bc $kb $k (: $prfarg $prms))))
    (: ($prfabs $prfarg) $ccln)))
```



```
;; Backward chainer with control (conditionals + context updaters)
(= (bc $kb (MkControl $absupd $argupd $bcont $rcont $mcont) $ctx (: $prf $ccln))
  (if ($bcont (: $prf $ccln) $ctx)
    (match $kb (: $prf $ccln) (if ($mcont (: $prf $ccln) $ctx) (: $prf $ccln) (empty)))
    (empty)))
(= (bc $kb (MkControl $absupd $argupd $bcont $rcont $mcont) $ctx (: ($prfabs $prfarg) $ccln))
  (if ($rcont (: ($prfabs $prfarg) $ccln) $ctx)
    (let* (((: $prfabs (-> $prms $ccln))
          (bc $kb (MkControl $absupd $argupd $bcont $rcont $mcont)
                ($absupd (: ($prfabs $prfarg) $ccln) $ctx) (: $prfabs (-> $prms $ccln))))
          ((: $prfarg $prms)
            (bc $kb (MkControl $absupd $argupd $bcont $rcont $mcont)
                  ($argupd (: ($prfabs $prfarg) $ccln) $ctx) (: $prfarg $prms))))
      (: ($prfabs $prfarg) $ccln))
    (empty)))
```

```
;; January precedes February, which precedes Mars, etc.
```

```
!(add-atom &kb (: JF (<= Jan Feb)))
```

```
!(add-atom &kb (: FM (<= Feb Mar)))
```

```
!(add-atom &kb (: MA (<= Mar Apr)))
```

```
!(add-atom &kb (: AM (<= Apr May)))
```

```
!(add-atom &kb (: MJ (<= May Jun)))
```

```
!(add-atom &kb (: JJ (<= Jun Jul)))
```

```
!(add-atom &kb (: JA (<= Jul Aug)))
```

```
!(add-atom &kb (: AS (<= Aug Sep)))
```

```
!(add-atom &kb (: SO (<= Sep Oct)))
```

```
!(add-atom &kb (: ON (<= Oct Nov)))
```

```
!(add-atom &kb (: ND (<= Nov Dec)))
```

```
;; Precedence is non strict, i.e. reflexive
```

```
!(add-atom &kb (: Refl (<= $x $x)))
```

```
;; Precedence is transitive
```

```
!(add-atom &kb (: Trans (-> (<= $x $y)
                             (-> (<= $y $z)
                                   (<= $x $z)))))
```

```
;; Shortcut rule: January precedes all months
```

```
!(add-atom &kb (: JPA (<= Jan $x)))
```

```

;; 1st observation: if
;; - the target theorem is (<= x x)
;; - the current proof is Refl
;; then continue.
!(add-atom &ctl-kb (: RS (Continue (: Refl $r) (MkTD (<= $x $x) $k))))

;; 2nd observation: if
;; - the target theorem is (<= Jan x)
;; - the current proof is JPA
;; then continue.
!(add-atom &ctl-kb (: JS (Continue (: JPA $r) (MkTD (<= Jan $x) $k))))

;; 3rd observation: if
;; - the target theorem is (<= $x $y) such that $x != Jan and $x != $y
;; - the current proof is Trans or FM to ND
;; then continue.
!(let $rn (superpose (Trans FM MA AM MJ JJ JA AS SO ON ND))
  (add-atom &ctl-kb (: TS (-> (!= Jan $x)
                              (-> (!= $x $y)
                                   (Continue (: $rn $rc)
                                             (MkTD (<= $x $y) $k))))))))

;; Backward chainer as continuation condition. Return True iff a
;; proof of continuation is found.
(: td-continuator (-> $a $ct Bool))
(= (td-continuator $query $ctx)
   (let $results (collapse (bc &ctl-kb &ctl (S (S 2)) (: $prf (Continue $query $ctx)))
                          (not (== () $results))))))

```

- Discover statistical patterns via reasoning.
- Formalize learning to reason about it.
- Use more universal control theory.
- Unify problem and control theory.

metta-lang.dev
github.com/trueagi-io/chaining