Towards Machine Learning Induction for Isabelle/HOL

This work was supported by the project AI&Reasoning (reg. no. CZ.02.1.01/0.0/0.0/15_003/0000466).



Yutaka Nagashima University of Innsbruck Czech Technical University



Yutaka Ng yutakang

Block or report user

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ITP (Inductive Theorem Proving) problems are at the heart of many verification and reasoning tasks in





Prof. Bernhard Gramlich https://www.logic.at/staff/gramlich/

Why induction?

Who is Isabelle?

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Why induction? we are convinced that substantial progress in ITP will take time.

spectacular breakthroughs are

<u>unrealistic</u>, in view of the enormous problems and the inherent difficulty of inductive theorem proving.

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git clone https://github.com/data61/PSL Interactive theorem proving with Isabelle/HOL







































 $\frac{\text{git clone https://github.com/data61/PSL}}{\text{lemma}} \text{ "map f (sep x xs) = sep (f x) (map f xs)"}$





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git clone https://github.com/data61/PSL

Try_Hard: the default strategy



prepa	ration	phase
picpa	auon	phase

How does PaMpeR work?

recommendation phase







STATISTICS

Archive of Formal Proofs (https://www.isa-afp.org)

	Statistics			
	Number of Articles: 46	8		
Home	Number of Authors: 313			
About	Lines of Code: \sim 2,170,300			
Submission				
Updating	Most used AFP article	Most used AFP articles:		
Entries	Name	Used by ? articles		
Using Entries	1. <u>Collections</u>	15		
Search	 <u>List-Index</u> <u>Coinductive</u> 	14 12		





















git clone https://github.com/data61/PSL

AITP2018 review



























proved theorem / subgoals / message



proof for the original goal, and auxiliary lemma <u>optimal</u> for proof automation

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Success story

PSL can find how to apply induction for easy problems.



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PSL can find how to apply induction for easy problems.

PaMpeR recommends which proof methods to use.



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CADE2017 ASE2018

Success story

PSL can find how to apply induction for easy problems.

PaMpeR recommends which proof methods to use.

PGT produces useful auxiliary lemmas.



CADE2017

ASE2018

Too good to be true?

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Too good to be true?

PSL can find how to apply induction for easy problema proof search only if PSL completes a proof search PaMpeR recommends which proof methods to use.



PGT produces useful auxiliary lemmas only if PSL with PGT completes a proof search

Too good to be true?

PSL can find how to apply induction for easy pretera proof search only if PSL completes a proof search **PaMpeR recommends which** proof methods to use but PaMpeR does not recommend arguments for proof methods **PGT produces useful auxiliary** lemmas only if PSL with PGT completes a proof search



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Recommend how to apply induction without completing a proof.

Too good to be true?

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Leave a star at GitHub for PSL!

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Let's write a review paper "AITP deserves High-Performance Computing, Too!"

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PaMpeR's feature extractor?











2017~ PaMpeR 2018 PaMpeR's data extraction 1986~ Isabelle Time 2004~ AFP 2018~ more articles in the AFP











Feature extractor?

lemma "map f (sep x xs) = sep (f x) (map f xs)"
Feature extractor? fun sep::"'a ⇒ 'a list ⇒ 'a list" where "sep a [] = []" | "sep a [x] = [x]" |

"sep a (x#y#zs) = x # a # sep a (y#zs)"

automatically proves and saves many auxiliary lemmas in the context sep.simps, sep.induct, sep.elims, etc.

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un-

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(map f (sep x xs) = sep (f x) (map f xs) remma map f (sep x xs) = find (find proof DInd)

At the time of development (2017), PSL does not know about

user defined constants (e.g. "sep") or user defined proof strategies (e.g. Dlnd). 2017: PSL define the "sep" function 2017: 2019: define the "DInd" strategy

Time