

# Project Proposal: Prediction by Compression

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Under reasonable conditions for  $C$ ,  $NCD_C$  approximates a metric

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Kolmogorov complexity  $K$ :

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$$NCD_K(s, t) = \frac{K(st) - \min(K(s), K(t))}{\max(K(s), K(t))}$$

$NCD_K$  is **the** distance metric:

$$\forall_{d,s,t} \text{ computable}(d) \Rightarrow NCD_K(s, t) \leq d(s, t)$$

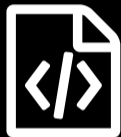


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Problem: Mathematical statements are short

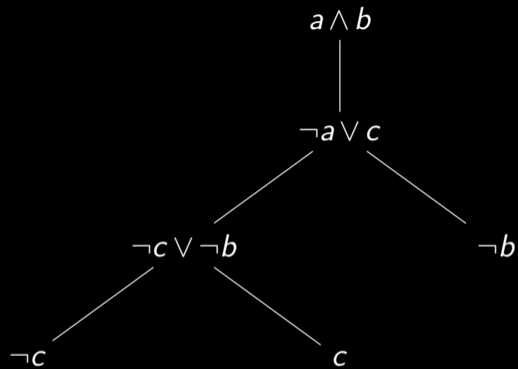
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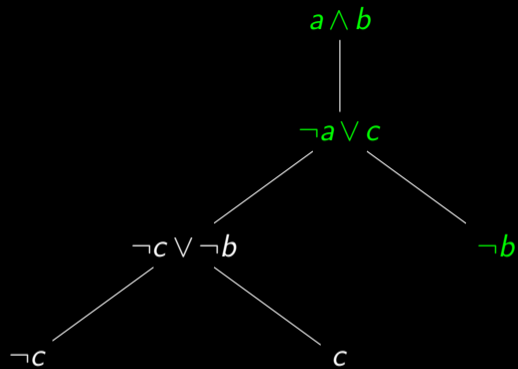
Compression: Prediction by Partial Matching

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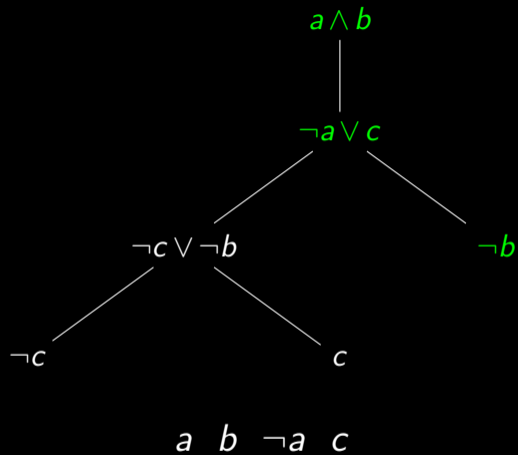
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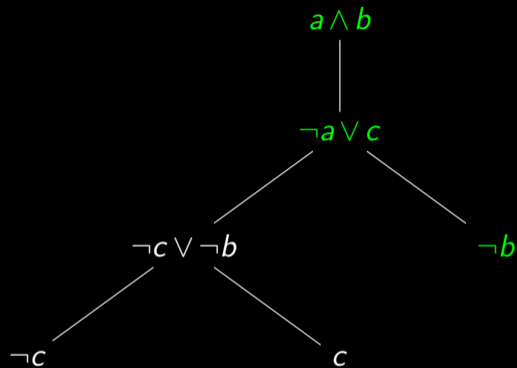
Compress entire proof states





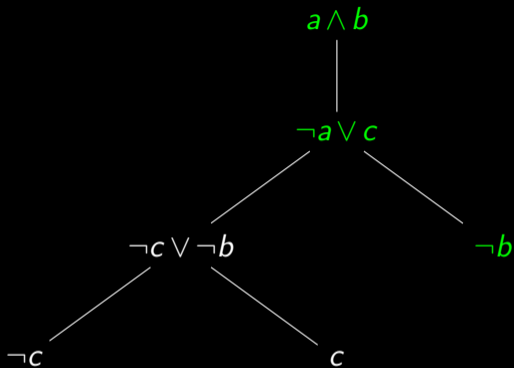






$a \quad b \quad \neg a \quad c$

“ $a \wedge b \rightarrow \neg a \vee c \rightarrow \neg b \wedge a \wedge b \rightarrow \neg a \wedge c$ ”

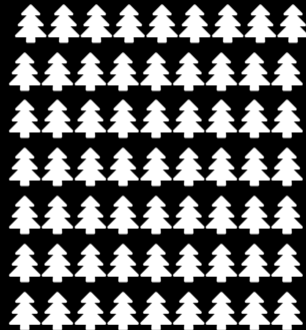


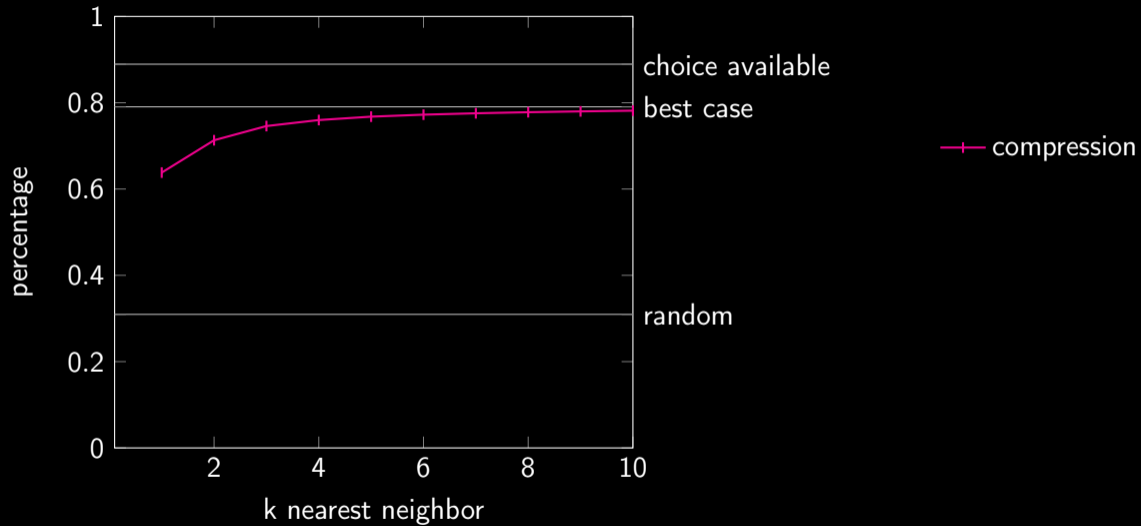
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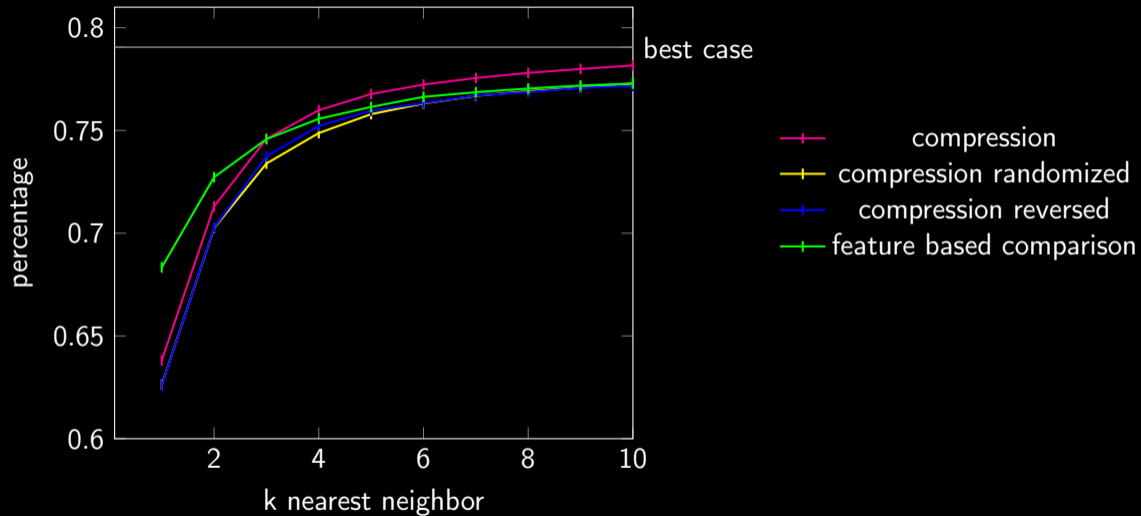
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Database







About 30-40 compressions per second

No vector space:  $n$  compressions per prediction

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Idea: Impose structure through an  $n$ -dimensional lattice

$$S_n = \{X \subseteq S \mid |X| = n\}$$

$$\text{out}(s) = \arg \max_{X \in S_n} \frac{\sum_{t, u \in X} \text{NCD}(t, u)}{\sum_{t \in X} \text{NCD}(s, t)}$$





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- ▷ No domain-specific knowledge required
- ▷ Predictions are competitive
- ▷ Robust against different representations of proof states

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- ▷ Impose a  $n$ -dimensional lattice on the data

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